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PATENT AND TRADEMARK OFFICEATTORNEY'S DOCKET NUMBER
2345/131**TRANSMITTAL LETTER TO THE UNITED STATES
DESIGNATED/ELECTED OFFICE (DO/EO/US)
CONCERNING A FILING UNDER 35 U.S.C. 371**

U.S. APPLICATION NO. (If known, see 37 CFR 1.5)

09/720616INTERNATIONAL APPLICATION NO
PCT/EP99/04438INTERNATIONAL FILING DATE
25 June 1999
(25.06.99)PRIORITY DATE CLAIMED.
26 June 1998
(26.06.98)

TITLE OF INVENTION

METHOD FOR CHECKING JAVA BYTE CODE PROGRAMMES FOR SECURITY CHARACTERISTICS

APPLICANT(S) FOR DO/EO/US

POSEGGA, Joachim; and VOGT, Harald

Applicant(s) herewith submits to the United States Designated/Elected Office (DO/EO/US) the following items and other information

1. ☒ This is a **FIRST** submission of items concerning a filing under 35 U.S.C. 371.
2. ☐ This is a **SECOND** or **SUBSEQUENT** submission of items concerning a filing under 35 U.S.C. 371.
3. ☒ This is an express request to begin national examination procedures (35 U.S.C. 371(f)) immediately rather than delay examination until the expiration of the applicable time limit set in 35 U.S.C. 371(b) and PCT Articles 22 and 39(1).
4. ☒ A proper Demand for International Preliminary Examination was made by the 19th month from the earliest claimed priority date.
5. ☒ A copy of the International Application as filed (35 U.S.C. 371(c)(2))
 - ☐ is transmitted herewith (required only if not transmitted by the International Bureau).
 - ☒ has been transmitted by the International Bureau.
 - ☐ is not required, as the application was filed in the United States Receiving Office (RO/US)
6. ☒ A translation of the International Application into English (35 U.S.C. 371(c)(2)).
7. ☒ Amendments to the claims of the International Application under PCT Article 19 (35 U.S.C. 371(c)(3))
 - ☐ are transmitted herewith (required only if not transmitted by the International Bureau).
 - ☐ have been transmitted by the International Bureau.
 - ☐ have not been made; however, the time limit for making such amendments has NOT expired.
 - ☒ have not been made and will not be made.
8. ☐ A translation of the amendments to the claims under PCT Article 19 (35 U.S.C. 371(c)(3)).
9. ☒ An oath or declaration of the inventor(s) (35 U.S.C. 371(c)(4)).
10. ☒ A translation of the annexes to the International Preliminary Examination Report under PCT Article 36 (35 U.S.C. 371(c)(5)).

Items 11. to 16. below concern other document(s) or information included:

11. An Information Disclosure Statement under 37 CFR 1.97 and 1.98.
12. ☒ An assignment document for recording. A separate cover sheet in compliance with 37 CFR 3.28 and 3.31 is included.
13. ☒ A **FIRST** preliminary amendment.
 - ☐ A **SECOND** or **SUBSEQUENT** preliminary amendment.
14. ☒ A substitute specification.
15. ☐ A change of power of attorney and/or address letter.
16. ☒ Other items or information: International Search Report and Preliminary Examination Report (translated).

Express Mail No.: EL594609247US

APPLICATION NO. (if known, see 37 CFR 1.1) 09/720616		INTERNATIONAL APPLICATION NO. PCT/EP99/04438		ATTORNEY'S DOCKET NUMBER 2345/131	
17. <input checked="" type="checkbox"/> The following fees are submitted: Basic National Fee (37 CFR 1.492(a)(1)-(5)): Search Report has been prepared by the EPO or JPO \$860.00 International preliminary examination fee paid to USPTO (37 CFR 1.482) . . . \$690.00 No international preliminary examination fee paid to USPTO (37 CFR 1.482) but international search fee paid to USPTO (37 CFR 1.445(a)(2)) \$710.00 Neither international preliminary examination fee (37 CFR 1.482) nor international search fee (37 CFR 1.445(a)(2)) paid to USPTO \$1,000.00 International preliminary examination fee paid to USPTO (37 CFR 1.482) and all claims satisfied provisions of PCT Article 33(2)-(4) \$100.00				CALCULATIONS PTO USE ONLY	
ENTER APPROPRIATE BASIC FEE AMOUNT =				\$ 860.00	
Surcharge of \$130.00 for furnishing the oath or declaration later than <input type="checkbox"/> 20 <input type="checkbox"/> 30 months from the earliest claimed priority date (37 CFR 1.492(e)).				\$	
Claims	Number Filed	Number Extra	Rate		
Total Claims	2 - 20 =	3	X \$18.00	\$ 0.00	
Independent Claims	1 - 3 =	0	X \$80.00	\$	
Multiple dependent claim(s) (if applicable)			+ \$270.00	\$	
TOTAL OF ABOVE CALCULATIONS =				\$ 860.00	
Reduction by 1/2 for filing by small entity, if applicable. Verified Small Entity statement must also be filed. (Note 37 CFR 1.9, 1.27, 1.28).				\$	
SUBTOTAL =				\$	
Processing fee of \$130.00 for furnishing the English translation later the <input type="checkbox"/> 20 <input type="checkbox"/> 30 months from the earliest claimed priority date (37 CFR 1.492(f)).				+	\$
TOTAL NATIONAL FEE =				\$ 860.00	
Fee for recording the enclosed assignment (37 CFR 1.21(h)). The assignment must be accompanied by an appropriate cover sheet (37 CFR 3.28, 3.31). \$40.00 per property				+	\$ 40.00
TOTAL FEES ENCLOSED =				\$ 900.00	
				Amount to be:	
				refunded	\$
				charged	\$

- a. ☐ A check in the amount of \$ _____ to cover the above fees is enclosed.
- b. ☒ Please charge my Deposit Account No. 11-0600 in the amount of \$900.00 to cover the above fees. A duplicate copy of this sheet is enclosed.
- c. ☒ The Commissioner is hereby authorized to charge any additional fees which may be required, or credit any overpayment to Deposit Account No. 11-0600. A duplicate copy of this sheet is enclosed.

NOTE: Where an appropriate time limit under 37 CFR 1.494 or 1.495 has not been met, a petition to revive (37 CFR 1.137(a) or (b)) must be filed and granted to restore the application to pending status

SEND ALL CORRESPONDENCE TO:

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SIGNATURE

Richard L. Mayer, Reg. No. 22,490

NAME

DATE

12/26/00

IN THE UNITED STATES PATENT AND TRADEMARK OFFICE

APPLICANT: POSEGGA et al.
SERIAL NO.: to be assigned
FILED: herewith
TITLE: METHOD FOR CHECKING JAVA BYTE CODE PROGRAMMES
FOR SECURITY CHARACTERISTICS
ART UNIT: not yet known
EXAMINER: not yet known

Assistant Commissioner for Patents
Washington, D.C. 20231

Sir:

PRELIMINARY AMENDMENT

Please amend the above-identified application before a first consideration on the merits as follows:

IN THE TITLE

Please amend the title to read as follows: --A METHOD FOR VERIFYING SAFETY PROPERTIES OF JAVA BYTE CODE PROGRAMS--.

IN THE FIGURES

Please add new Fig. 1, submitted herewith.

IN THE SPECIFICATION

Please replace the specification of record, i.e., the English translation of the PCT application text (submitted herewith) with the substitute specification submitted herewith. The substitute specification is being submitted due to the nature and number of amendments.

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A marked-up copy of the substitute specification is also submitted herewith, in accordance with 37 C.F.R. 1.125(b). It is respectfully submitted that the substitute specification contains no new matter.

On page 15, line 1, of the specification of record change "Patent Claim" to --WHAT IS CLAIMED IS:--.

IN THE CLAIMS

Please cancel without prejudice claim 1 and add new claims 2-3 as follows:

--2. (new) A method for verifying safety properties of a Java byte code program, the method comprising:

mapping a functioning of the byte code program by a potentially infinite state transition system onto a finite state transition system using an algorithm describing first properties of byte code instructions,

mapping a state space of an interpreter onto a finite set of states in the finite state transition system, information not needed for a checking of an acceptability of the byte code program being omitted, so that the finite state transition system contains only type information useable for the checking of the acceptability of the byte code program;

entering the type information useable for the checking of the acceptability of the byte code program into a model checker;

determining second properties which characterize an acceptable byte code program using a logic operation including formulas;

entering the determined second properties which characterize an acceptable byte code program as conditional set into the model checker, the conditional set including a plurality of individual conditions;

interpreting, using the model checker, each of the plurality of individual conditions as a specification language for system properties of the byte code program;

verifying, using the model checker, whether each of the plurality of individual conditions is fulfilled by the state transition system; and then

automatically releasing the byte code program for further processing when the state transition system fulfills all of the plurality of individual conditions.

3. (new) The method as recited in claim 2 wherein the interpreter is a java virtual machine.--.

IN THE ABSTRACT

Please delete the abstract of record and replace it with the following amended abstract:
--In a method for verifying the safety properties of Java byte code programs, the functioning of the byte code program to be verified is modeled on a finite state transition system M, and the state space of the Java Virtual Machine (JVM) on a finite set of states in M. Once entered into a model checker, the data of finite state transition system M are compared to the data entered in the model checker as conditional set S to determine properties characterizing an acceptable byte code program. The byte code program to be checked is released for further processing only when the state transition system M fulfills all conditions of conditional set S.--.

REMARKS

This Preliminary Amendment cancels original claim 1 and adds new claims 2-3, and adds new Fig. 1. The new claims and figure do not add new matter to the application but do conform the claims to U.S. Patent and Trademark Office rules.

The amendments to the specification and abstract are to conform the specification and abstract to U.S. Patent and Trademark Office rules. It is respectfully submitted that the amendments to the specification and abstract do not introduce new matter into the application.

The underlying PCT application No. PCT/EP99/04438 includes a Search Report, a copy of which is included herewith.

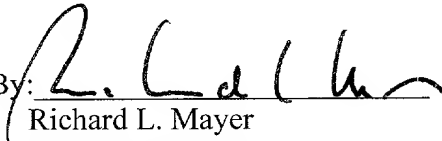
Conclusion

Consideration of the present application as amended is hereby respectfully requested.

Respectfully Submitted,

Kenyon & Kenyon

Dated: 12/26/00

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[2345/131]

A METHOD FOR ~~{CHECKING}~~ VERIFYING SAFETY PROPERTIES OF JAVA
BYTE CODE ~~{PROGRAMMES FOR SECURITY CHARACTERISTICS}~~
PROGRAMS

~~{Description:}~~ Field of the Invention

The present invention relates to a method for verifying the safety properties of Java byte code programs ~~[in accordance with Claim 1.]~~ using byte code verification.

5 Related Technology

Java is a programming language developed by Sun Microsystems that is translated into so-called byte code programs. Although originally developed for Java, byte code is suited as a target language for other programming languages, and compilers (such as ADA) are already available for this purpose.

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Java's unique concept is to store programs in a network where network subscribers can retrieve them and run them on one of their units. This enables, for example, units to be configured and serviced from remote locations. However, it is also interesting in the context of smart cards, since it can be used to enhance their functionality. The

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following properties are characteristic of byte code programs and are significant for their intended application:

- Byte code programs are compact. Byte code instructions offer substantially greater functionality than, for example, computer language instructions. Concepts, such as object orientation, are modeled on Java to enable them to be directly supported.
- Byte code programs can be interpreted efficiently, which makes them independent from a specific destination machine. The only requirement is an interpreter (the so-called "Java virtual machine", JVM) on the destination machine, to be able to run any desired byte code programs. The interpreter

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itself can be implemented in a compact form and be integrated in (virtually) every unit.

5 The JVM interprets a character string input as a byte code program. In the process, the JVM does not verify whether the character string is actually a byte code program. It processes the characters as byte code commands, without performing a more precise test to check for their correctness. The JVM does not verify the byte codes, since this would significantly degrade the execution speed. In principle, therefore, it is possible to interpret any string of characters as a byte code program. Thus, - given a more
10 precise knowledge of the JVM implementation - one could even circumvent its safety mechanisms and spy out data on the computer running the program. On the other hand, because of the language design, a secure byte code program can only access data and resources of the destination computer by calling upon functions of the JVM for that purpose.

15 To prevent character strings from being processed, which do not constitute permissible (safe) byte code programs, a process, the so-called byte code verification, is added. It tests the character string to see if it conforms with certain requirements. These requirements ensure that a character sequence is a safe byte code program. Since this
20 verification only needs to be done once, the run time is not adversely affected with respect to the execution speed. The byte code verification process tests a string of characters to check whether it constitutes a permissible byte code program and, thus, poses no danger to the integrity of the unit running the program and can, thus, be transferred to the JVM for interpretation (see T. Lindholm, F. Yellin, "The Java Virtual
25 Machine Specification"; Sun Microsystems; Addison-Wesley, 1996, which is hereby incorporated by reference herein).^[1]

A permissible byte code program is characterized by certain properties, which are described as "structural constraints" in T. Lindholm, F. Yellin: "The Java Virtual
30 Machine Specification", Sun Microsystems, 1996, which is hereby incorporated by

reference herein. For the most part, they describe the safety, with respect to type, of a byte code program. This means that, in a byte code program, the instructions are arranged in such a way that an instruction always receives those data as parameters that correspond, in type, to that which is expected of the instruction. These properties ensure that the byte code program remains under the control of the JVM and cannot directly access computer resources.

The byte code verification process is described by two actualities: first by the description of the properties that it is supposed to guarantee for a byte code program; secondly by the implementation of the process (in this case, a C program). Neither of the two descriptions is suited for making formal, machine calculations, and, therefore, for formally guaranteeing the safety of byte code programs.

Also known is a tool, referred to as a model checker, for verifying state transition systems (see K. L. McMillan "Symbolic Model Checking" Kluwer Academic Publishers, 1993, which is hereby incorporated by reference herein).~~f).]~~

State transition systems make up a general model for programs that run on computers or other machines (such as the JVM). In this context, ~~[one takes into consideration]~~ the states that the machine assumes during execution of the program and which transitions, i.e., state changes are possible.

Model checking is a technique which can only be used to verify finite state transition systems. For this purpose, a model checker examines the entire state space of the system to calculate the properties of the system. A typical application of model checkers is verifying whether a system satisfies certain properties. For this, as an input, the model checker requires a description of the system to be tested, as well as a description of the properties whose validity needs to be established in the system. Both cases have to do with formal descriptions, i.e., descriptions whose formal semantics is established, and on which mathematical calculations can, therefore, be

performed.

Generally, software systems, including byte code programs, have an infinite state space. This means that the model checking method cannot be applied to byte code programs characterized by an infinite state space.

Summary of the Invention

An object of the present invention is to provide ~~[The technical objective is directed to]~~ a method which will ensure a ~~[greatest possible]~~ high degree of security in verifying the safety properties of byte code programs.

The method according to the present invention ~~[starts out from the factual situation that]~~ may be applied to a byte code program ~~[contains]~~ containing a string of characters (bytes), each character either being interpreted as an instruction or as data (as an input for the preceding instruction). Thus, a byte code program represents a series of instructions and corresponding data. Therefore, it can also be interpreted as a description of a state transition system that transforms states of the JVM. The states of the system are identified on the basis of the states of the JVM, and the transitions are defined by the program instructions.

In accordance with the present invention, suitable mapping is used to map the state transition system described above (having a potentially infinite state space) onto a system having a finite number of states. A finite number of states is abstracted in that the system rejects all information not needed in determining whether the original byte code program is valid. This is ~~[essentially]~~ accomplished by replacing the specific, or concrete, data values with pure, or abstract, type information. This can be done because the validity of a byte code program does not depend on specific values. In the process, however, all information needed to check the validity of the byte code program is retained. The resulting state transition system M has a finite number of states and thus fulfills the basic condition for processing in a model checker. The

resulting state transition system M is entered into the model checker.

In a further step, the properties which characterize a valid byte code program, as T. Lindholm, F. Yellin describe in "The Java Virtual Machine Specification", Sun Microsystems, 1996 as "structural constraints", are formulated in a logic understood by the model checker as a specification language for the system properties. The result is a set of formulas that are transferred to the model checker as a conditional set S, to be used as a basis of comparison for state transition system M. The above procedure ensures that the formulas only relate to information that is present both in the original, potentially infinite-state system, as well as in the finite-state system derived therefrom. The model checker is started with the inputs just described. As a result, it provides information as to whether state transition system M fulfills system properties described in the specification language as conditional set S in the form of formulas. In the process, the model checker verifies whether each condition s of conditional set S is fulfilled by state transition system M of the byte code program to be checked. If each condition is fulfilled, then it is ensured that the original byte code program is a valid program that does not threaten the safety of the computer executing the program. The byte code program is released for further computer processing.

However, if the data of the byte code program to be checked acquired in state transition system M do not fulfill the system properties described in conditional set S in the form of formulas and characterizing a byte code program, then ~~[the fundamental assumption is]~~ it may be assumed that the tested byte code program can threaten the safety of the computer. The tested byte code program is not released for further processing.

Brief Description of the Drawings

The method according to the present invention is ~~[elucidated in the following:]~~ elaborated upon in the following with reference to the drawing, in which:
~~[The purpose of the]~~ Fig. 1 shows a flow chart of a method for verifying the safety properties of Java byte code programs.

Detailed Description

Referring to Fig. 1, the method according to the present invention maps a functioning of the byte code program by a potentially infinite state transition system onto a finite state transition system using an algorithm describing first properties of byte code instructions (block 102). A state space of an interpreter is mapped onto a finite set of states in the finite state transition system, information not needed for a checking of an acceptability of the byte code program being omitted, so that the finite state transition system contains only type information useable for the checking of the acceptability of the byte code program (block 104). The type information useable for the checking of the acceptability of the byte code program is entered into a model checker (block 106). Second properties which characterize an acceptable byte code program are determined using a logic operation including formulas (block 108). The determined second properties which characterize an acceptable byte code program are entered as conditional set into the model checker, the conditional set including a plurality of individual conditions (block 110). Using the model checker, each of the plurality of individual conditions is interpreted as a specification language for system properties of the byte code program (block 112). Using the model checker, whether each of the plurality of individual conditions is fulfilled by the state transition system is verified (block 114). Then, the byte code program is automatically released for further processing when the state transition system fulfills all of the plurality of individual conditions (see block 116).

A purpose of a "byte code verification" is to check a class data file, thus the unit in which a Java class description is summarized, to verify that it can be safely executed. For this, the individual methods of the class (referred to here as byte code programs) are individually released and subjected to the actual byte code verification.

A class data file contains additional data, ~~[essentially]~~ basically constants, which are referred to in the byte code programs. For the subsequent steps, it is advantageous to resolve, or sort out, these references in the byte code programs and incorporate them in the programs before beginning with the actual verification process.

This preprocessing provides one with a description of the byte code programs, which does not differ substantially from the original form.

The JVM, the standard interpreter for byte code programs, is described by an abstract machine, i.e., by a set of states. The execution of a byte code program corresponds to the interpretation of byte code instructions by, or on the basis of, state transitions. In this context, a byte code instruction effectively transforms the active state of the JVM into a new state.

Thus, a byte code program defines a state transition system, the state space being set by the JVM and the transitions being defined by the instructions of the byte code program. This state transition system has a potentially infinite state space. Therefore, it is not suited for verification by a model checker. For that, it ~~has to be~~ is mapped onto a finite state transition system M. This mapping process is described in the following.

One does not first construct an infinite transition system that is subsequently mapped onto a finite transition system. Rather, a finite state transition system M is constructed directly from the byte code program, using rules which describe the properties of the byte code instructions.

The ~~method of~~ functioning of a byte code program is mapped onto a finite state transition system M. The state space of the JVM is mapped onto a finite set of states in state transition system M. The byte code instructions then effect transitions on this finite state quantity, its principal effect being retained. State transition system M is described in a way that allows the description to be used as an input for the model checker.

The input for the model checker has two parts:

1. A system description, designated as state transition system M, including
 - a description of the state space, defined by the state variables and their domains;

- a provision specifying the start states of the system;
- a transition relation indicating how states are transformed and under which conditions.

2. A set of specifications designated as conditional set S, describing the desired properties of a permissible byte code program. These properties are required by state transition system M, which is described by the byte code program. Thus, safety can be ensured when executing the byte code program. These properties are valid in the specific, infinite state transition system (and thus in the original byte code program) when they are valid in the abstract, finite state transition system M. The model checker verifies whether they are valid in the abstract state transition system M. If they are valid, then ~~one~~ it can ~~infer~~ be inferred that the properties are also fulfilled by specific system and, thus, in the original program.

15 The following describes the generation of the required data:

The state space of state transition system M is constructed of the following components:

- an instruction counter;
- an operand stack, which records the parameters and results of instructions;
- 20 - a field of local variables;
- a field of global variables;
- variables for keeping bookkeeping information; one can refer to this to describe system properties which cannot be described by the remaining components. This includes, inter alia: ~~a~~
- 25 - a stack, which contains information on the active subroutines (subroutines in a byte code program).

The values of variables and stack elements are derived from finite domains. The instruction counter can assume a finite number of address values. The operand stack is limited in magnitude. Overall, therefore, the state space of M is finite.

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The instruction directory V contains descriptions of all byte code instructions,

including the following information for each instruction:

- the description of the instruction (1 byte)
- the number of parameters in bytes (characters)
- a description of the effect on states of the JVM
- 5 • conditions, which must be met by a state of the JVM, so that the instruction can be applied (C1)
- conditions (C2) which must be met by the JVM because the instruction occurs in the program. These conditions do not necessarily relate to individual states, but also to execution paths, i.e., a series of states.

10 For the conditions in V, the following property (A1) applies: the description of the conditions in instruction directory V uses only those symbols used in the description of the state space for M.

15 From a byte code program, a compiler generates a finite state transition system. The byte code program exists as a string of characters in the input, it being possible for an individual character to be interpreted as an instruction or parameter of an instruction. The first character is interpreted as an instruction.

20 The goal of the compilation is to construct a state transition system M and a conditional set S. The compilation process is described as follows:

1. ~~{One begins}~~ Start with an initial ("empty") transition system M. The state space of M is already defined by the abstract representation of the JVM state space. However, the start state of the JVM defines the start state of M. However, no transitions between the states are defined yet.
- 25 2. ~~{One begins}~~ Start with an empty conditional set S.
3. As long as characters are present in the input, ~~{one implements}~~ implement the following steps 4 through 8:
4. ~~{One reads}~~ Read a character from the input; this describes program instruction B.
- 30 5. ~~{One looks}~~ Look up in the instruction directory V to check which parameters B requires.

6. Read the parameters P required for B from the input.
7. Transfer (B,P) to the abstraction component.
8. Transfer (B,P) to the conditional component.

5 From an instruction B and corresponding parameters P, the abstraction component produces a set of state transitions. State transition system M is expanded by these transitions.

1. ~~{One looks}~~ Look up in the instruction directory V to determine what effect B has on a state of the JVM. The effect is described by a rule so that the follow-up state of the JVM results from a calculation made on the active state. (An explicit indication of the possible transitions is not possible since the JVM has a (virtually) infinite state space.)
- 10 2. From this, ~~{one derives}~~ derive a description of the effect on states of M under consideration of P. The parameters P enter into the calculation of the follow-up state. If ~~{one applies}~~ a JVM rule is applied to an M state, then ~~{one obtains}~~ a set of M states which represent the possible follow-up states of the JVM is obtained. This yields a set of follow-up states for M since it is not the specific values in P, but rather only the information relevant for M that is considered. Which JVM follow-up state is actually assumed when the byte code program is executed depends on the specific values in P.
- 15 3. This description defines a set number of state transitions; M expanded by these transitions. These transitions can be described by a rule or be explicitly specified. The latter is possible since M only knows a finite number of states and, therefore, all transitions can be specified by state pairs. Ultimately, the description follows that which is understood by the model checker being used.
- 20 For practical reasons, one explicitly avoids specifying all transitions since the data set for this is substantially greater than in the case of implicit specification on the basis of rules.
- 25
- 30

From (B,P), the conditional component produces a set of conditions which are ultimately to be verified by the model checker. The conditions can be

represented in different ways, for instance by time-logic formulas.

1. ~~{One looks}~~ Look up in the instruction directory V to determine which conditions C1 are stipulated for the active JVM state so that B can be executed.
2. ~~{One transfers}~~ Transfer C1 to corresponding M-conditions. In accordance with (A1), the conditions in V and the state description for M are adjusted to one another in such a way that C1 can be interpreted for M-states. The transfer to M essentially includes using P to produce specific instances of C1, and translating them into a logic formula.
3. ~~{One looks}~~ Look up accordingly the conditions C2 which arise for the entire system M through the use of B.
4. ~~{One transfers}~~ Transfer accordingly C2 to M.
5. ~~{One expands}~~ Expand the conditional set S by the representations for C1 and C2.

The described construction of state transition system M and of conditional set S for a byte code program to be verified proceeds ~~[fully]~~ automatically. The instruction directory V and its components are established one time, in the same way as the state space of M, for the specific purpose of the “byte code verification”, i.e., of the verification of safety properties of byte code programs.

The method of operation of the model checker is explained in the following. For each condition s from conditional set S, the model checker verifies if it is fulfilled by state transition system M. If a condition s is not fulfilled, the model checker generates information that can be used together with the knowledge of the instruction for which s was produced, to analyze how the conditions for a secure execution of the JVM can be violated by the byte code program. A typical output of the model checker is made up of the indication of an execution path, i.e., of a sequence of M-states, which leads to a violation of certain conditions.

A precise analysis of the execution path is not needed. The aim of the byte code

verification is merely to establish whether or not the program can violate the conditions. Information pertaining to how a violation can take place may be interesting, but is not needed, since it is merely necessary to reject potentially unsafe programs. In the process, one has to accept the fact that safe programs will also be rejected, which do, in fact, violate conditions, but, when more thoroughly analyzed, would reveal that in an actual environment, they would not be able to cause damage. Since programs of this kind are not generally developed by compilers, but rather only by deliberate encoding, significant restrictions are not being made in this case.

Controlling the byte code verification lies in coordinating the various components. A class data set is subjected to the byte code verification, in that

1. the byte code programs (methods) are released individually;
2. the references in the byte code program are resolved (pre-processing);
3. a state system and a conditional set are constructed;
4. the model checker is started with this input;
5. given success of the model checker, it is repeated for the next byte code program; in the case of failure, it is reported that the byte code verification for the class data set being analyzed has failed.

The method in accordance with the present invention ~~[can still]~~ may be expanded by a few points. Which properties can be verified depends on the manner in which a byte code program is mapped onto a finite state system. To establish, for instance, the acceptability of a byte code program, ~~[one can use]~~ the mapping outlined above may be used. By selecting another abstraction mapping, ~~[one can verify]~~ other properties ~~[An interesting]~~ may be verified. An useful property would be, for instance, restricting to a certain extent the consumption of resources by a byte code program. Through the method of the present invention, the byte code verification concept that is critical to safety can be implemented on a formal basis. In this manner, ~~[the greatest possible]~~ a high level of safety [is] may be attained for this aspect of Java technology. ~~[One can expect]~~ It may be expected that, above and beyond that, the technique can also be used to verify more far-reaching properties of Applets, making a larger range of applications

interesting.

This technique ~~[is particularly interesting]~~ may be applied in the environment of Java-capable smart cards (Java-capable meaning the ability to execute byte code programs) where the safety requirements are especially stringent. Java-capable smart cards make it possible to load and execute programs and, in fact, even when they are already in the possession of the end customer (this is not possible or only possible to a limited extent when conventional smart cards are used). In this context, it is ~~[extremely]~~ important that only those programs are loaded and executed which do not violate the integrity of the card and the data stored thereon, since misusing such data can cause considerable personal or financial damage.

The safety of byte code programs ~~[can]~~ may be guaranteed by using the described technique, and a certain functionality can be assured through additional upgrades. This ~~[gives one]~~ provides greater confidence in applications designed to be run on platforms that are critical to safety, such as smart cards. [

Reference Symbol List

JVM interpreter for Java byte code programs (Java virtual machine)

~~M~~ finite state transition system

~~S~~ conditional set

~~s~~ condition of S

~~V~~ instruction directory for byte code instructions

~~C1~~ conditions that have to be met by a state of the JVM

$C2$ conditions that have to be met by the JVM

$A1$ property that applies for conditions in V

5 B program instruction

P parameters for B

00000000000000000000000000000000

[2345/131]

METHOD FOR CHECKING JAVA BYTE CODE PROGRAMMES FOR
SECURITY CHARACTERISTICS

Description:

The present invention relates to a method for verifying the safety properties of Java byte code programs in accordance with Claim 1.

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Java is a programming language developed by Sun Microsystems that is translated into so-called byte code programs. Although originally developed for Java, byte code is suited as a target language for other programming languages, and compilers (such as ADA) are already available for this purpose.

10

Java's unique concept is to store programs in a network where network subscribers can retrieve them and run them on one of their units. This enables, for example, units to be configured and serviced from remote locations. However, it is also interesting in the context of smart cards, since it can be used to enhance their functionality. The following properties are characteristic of byte code programs and are significant for their intended application:

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- Byte code programs are compact. Byte code instructions offer substantially greater functionality than, for example, computer language instructions. Concepts, such as object orientation, are modeled on Java to enable them to be directly supported.
- Byte code programs can be interpreted efficiently, which makes them independent from a specific destination machine. The only requirement is an interpreter (the so-called "Java virtual machine", JVM) on the destination machine, to be able to run any desired byte code programs. The interpreter itself can be implemented in a compact form and be integrated in (virtually)

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every unit.

The JVM interprets a character string input as a byte code program. In the process, the JVM does not verify whether the character string is actually a byte code program. It processes the characters as byte code commands, without performing a more precise test to check for their correctness. The JVM does not verify the byte codes, since this would significantly degrade the execution speed. In principle, therefore, it is possible to interpret any string of characters as a byte code program. Thus, - given a more precise knowledge of the JVM implementation - one could even circumvent its safety mechanisms and spy out data on the computer running the program. On the other hand, because of the language design, a secure byte code program can only access data and resources of the destination computer by calling upon functions of the JVM for that purpose.

To prevent character strings from being processed, which do not constitute permissible (safe) byte code programs, a process, the so-called byte code verification, is added. It tests the character string to see if it conforms with certain requirements. These requirements ensure that a character sequence is a safe byte code program. Since this verification only needs to be done once, the run time is not adversely affected with respect to the execution speed. The byte code verification process tests a string of characters to check whether it constitutes a permissible byte code program and, thus, poses no danger to the integrity of the unit running the program and can, thus, be transferred to the JVM for interpretation (see T. Lindholm, F. Yellin, "The Java Virtual Machine Specification"; Sun Microsystems; Addison-Wesley, 1996).

A permissible byte code program is characterized by certain properties, which are described as "structural constraints" in T. Lindholm, F. Yellin: "The Java Virtual Machine Specification", Sun Microsystems, 1996. For the most part, they describe the safety, with respect to type, of a byte code program. This means that, in a byte code program, the instructions are arranged in such a way that an instruction always receives

those data as parameters that correspond, in type, to that which is expected of the instruction. These properties ensure that the byte code program remains under the control of the JVM and cannot directly access computer resources.

5 The byte code verification process is described by two actualities: first by the description of the properties that it is supposed to guarantee for a byte code program; secondly by the implementation of the process (in this case, a C program). Neither of the two descriptions is suited for making formal, machine calculations, and, therefore, for formally guaranteeing the safety of byte code programs.

10 Also known is a tool, referred to as a model checker, for verifying state transition systems (see K. L. McMillan "Symbolic Model Checking" Kluwer Academic Publishers, 1993).

15 State transition systems make up a general model for programs that run on computers or other machines (such as the JVM). In this context, one takes into consideration the states that the machine assumes during execution of the program and which transitions, i.e., state changes are possible.

20 Model checking is a technique which can only be used to verify finite state transition systems. For this purpose, a model checker examines the entire state space of the system to calculate the properties of the system. A typical application of model checkers is verifying whether a system satisfies certain properties. For this, as an input, the model checker requires a description of the system to be tested, as well as a
25 description of the properties whose validity needs to be established in the system. Both cases have to do with formal descriptions, i.e., descriptions whose formal semantics is established, and on which mathematical calculations can, therefore, be performed.

30 Generally, software systems, including byte code programs, have an infinite state

space. This means that the model checking method cannot be applied to byte code programs characterized by an infinite state space.

5 The technical objective is directed to a method which will ensure a greatest possible security in verifying the safety properties of byte code programs.

The method according to the present invention starts out from the factual situation that a byte code program contains a string of characters (bytes), each character either being interpreted as an instruction or as data (as an input for the preceding instruction).

10 Thus, a byte code program represents a series of instructions and corresponding data. Therefore, it can also be interpreted as a description of a state transition system that transforms states of the JVM. The states of the system are identified on the basis of the states of the JVM, and the transitions are defined by the program instructions.

15 In accordance with the present invention, suitable mapping is used to map the state transition system described above (having a potentially infinite state space) onto a system having a finite number of states. A finite number of states is abstracted in that the system rejects all information not needed in determining whether the original byte code program is valid. This is essentially accomplished by replacing the concrete data values with abstract type information. This can be done because the validity of a byte
20 code program does not depend on specific values. In the process, however, all information needed to check the validity of the byte code program is retained. The resulting state transition system M has a finite number of states and thus fulfills the basic condition for processing in a model checker. The resulting state transition
25 system M is entered into the model checker.

In a further step, the properties which characterize a valid byte code program, as T. Lindholm, F. Yellin describe in "The Java Virtual Machine Specification", Sun
30 Microsystems, 1996 as "structural constraints", are formulated in a logic understood by the model checker as a specification language for the system properties. The result is a

set of formulas that are transferred to the model checker as a conditional set S, to be used as a basis of comparison for state transition system M. The above procedure ensures that the formulas only relate to information that is present both in the original, potentially infinite-state system, as well as in the finite-state system derived therefrom.

5

The model checker is started with the inputs just described. As a result, it provides information as to whether state transition system M fulfills system properties described in the specification language as conditional set S in the form of formulas. In the process, the model checker verifies whether each condition s of conditional set S is fulfilled by state transition system M of the byte code program to be checked. If each condition is fulfilled, then it is ensured that the original byte code program is a valid program that does not threaten the safety of the computer executing the program. The byte code program is released for further computer processing.

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However, if the data of the byte code program to be checked acquired in state transition system M do not fulfill the system properties described in conditional set S in the form of formulas and characterizing a byte code program, then the fundamental assumption is that the tested byte code program can threaten the safety of the computer. The tested byte code program is not released for further processing.

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The method according to the present invention is elucidated in the following:
The purpose of the "byte code verification" is to check a class data file, thus the unit in which a Java class description is summarized, to verify that it can be safely executed. For this, the individual methods of the class (referred to here as byte code programs) are individually released and subjected to the actual byte code verification.
A class data file contains additional data, essentially constants, which are referred to in the byte code programs. For the subsequent steps, it is advantageous to resolve these references in the byte code programs and incorporate them in the programs before beginning with the actual verification process.

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This preprocessing provides one with a description of the byte code programs, which

does not differ substantially from the original form.

The JVM, the standard interpreter for byte code programs, is described by an abstract machine, i.e., by a set of states. The execution of a byte code program corresponds to the interpretation of byte code instructions on the basis of state transitions. In this context, a byte code instruction effectively transforms the active state of the JVM into a new state.

Thus, a byte code program defines a state transition system, the state space being set by the JVM and the transitions being defined by the instructions of the byte code program. This state transition system has a potentially infinite state space. Therefore, it is not suited for verification by a model checker. For that, it has to be mapped onto a finite state transition system M. This mapping process is described in the following.

One does not first construct an infinite transition system that is subsequently mapped onto a finite transition system. Rather, a finite state transition system M is constructed directly from the byte code program, using rules which describe the properties of the byte code instructions.

The method of functioning of a byte code program is mapped onto a finite state transition system M. The state space of the JVM is mapped onto a finite set of states in state transition system M. The byte code instructions then effect transitions on this finite state quantity, its principal effect being retained. State transition system M is described in a way that allows the description to be used as an input for the model checker.

The input for the model checker has two parts:

1. A system description, designated as state transition system M, including
 - a description of the state space, defined by the state variables and their domains;
 - 5 - a provision specifying the start states of the system;
 - a transition relation indicating how states are transformed and under which conditions.
2. A set of specifications designated as conditional set S, describing the desired properties of a permissible byte code program. These properties are required by state transition system M, which is described by the byte code program.
10 Thus, safety can be ensured when executing the byte code program. These properties are valid in the specific, infinite state transition system (and thus in the original byte code program) when they are valid in the abstract, finite state transition system M. The model checker verifies whether they are valid in the
15 abstract state transition system M. If they are valid, then one can infer that the properties are also fulfilled by specific system and, thus, in the original program.

The following describes the generation of the required data:

20 The state space of state transition system M is constructed of the following components:

- an instruction counter;
 - an operand stack, which records the parameters and results of instructions;
 - a field of local variables;
 - 25 - a field of global variables;
 - variables for keeping bookkeeping information; one can refer to this to describe system properties which cannot be described by the remaining components.
- This includes, inter alia: a
- stack, which contains information on the active subroutines (subroutines in a
30 byte code program).

The values of variables and stack elements are derived from finite domains. The

instruction counter can assume a finite number of address values. The operand stack is limited in magnitude. Overall, therefore, the state space of M is finite.

The instruction directory V contains descriptions of all byte code instructions, including the following information for each instruction:

- the description of the instruction (1 byte)
- the number of parameters in bytes (characters)
- a description of the effect on states of the JVM
- conditions, which must be met by a state of the JVM, so that the instruction can be applied (C1)
- conditions (C2) which must be met by the JVM because the instruction occurs in the program. These conditions do not necessarily relate to individual states, but also to execution paths, i.e., a series of states.

For the conditions in V, the following property (A1) applies: the description of the conditions in instruction directory V uses only those symbols used in the description of the state space for M.

From a byte code program, a compiler generates a finite state transition system. The byte code program exists as a string of characters in the input, it being possible for an individual character to be interpreted as an instruction or parameter of an instruction. The first character is interpreted as an instruction.

The goal of the compilation is to construct a state transition system M and a conditional set S. The compilation process is described as follows:

1. One begins with an initial (“empty”) transition system M. The state space of M is already defined by the abstract representation of the JVM state space. However, the start state of the JVM defines the start state of M. However, no transitions between the states are defined yet.
2. One begins with an empty conditional set S.
3. As long as characters are present in the input, one implements the following steps 4 through 8:

4. One reads a character from the input; this describes program instruction B.
5. One looks up in the instruction directory V to check which parameters B requires.
6. Read the parameters P required for B from the input.
- 5 7. Transfer (B,P) to the abstraction component.
8. Transfer (B,P) to the conditional component.

10 From an instruction B and corresponding parameters P, the abstraction component produces a set of state transitions. State transition system M is expanded by these transitions.

1. One looks up in the instruction directory V to determine what effect B has on a state of the JVM. The effect is described by a rule so that the follow-up state of the JVM results from a calculation made on the active state. (An explicit indication of the possible transitions is not possible since the JVM has a
15 (virtually) infinite state space.)
2. From this, one derives a description of the effect on states of M under consideration of P. The parameters P enter into the calculation of the follow-up state. If one applies a JVM rule to an M state, then one obtains a set of M states which represent the possible follow-up states of the JVM. This yields a
20 set of follow-up states for M since it is not the specific values in P, but rather only the information relevant for M that is considered. Which JVM follow-up state is actually assumed when the byte code program is executed depends on the specific values in P.
3. This description defines a set number of state transitions; M expanded by these
25 transitions. These transitions can be described by a rule or be explicitly specified. The latter is possible since M only knows a finite number of states and, therefore, all transitions can be specified by state pairs. Ultimately, the description follows that which is understood by the model checker being used. For practical reasons, one explicitly avoids specifying all transitions since the
30 data set for this is substantially greater than in the case of implicit specification

on the basis of rules.

From (B,P), the conditional component produces a set of conditions which are ultimately to be verified by the model checker. The conditions can be represented in different ways, for instance by time-logic formulas.

1. One looks up in the instruction directory V to determine which conditions C1 are stipulated for the active JVM state so that B can be executed.
2. One transfers C1 to corresponding M-conditions. In accordance with (A1), the conditions in V and the state description for M are adjusted to one another in such a way that C1 can be interpreted for M-states. The transfer to M essentially includes using P to produce specific instances of C1, and translating them into a logic formula.
3. One looks up accordingly the conditions C2 which arise for the entire system M through the use of B.
4. One transfers accordingly C2 to M.
5. One expands the conditional set S by the representations for C1 and C2.

The described construction of state transition system M and of conditional set S for a byte code program to be verified proceeds fully automatically. The instruction directory V and its components are established one time, in the same way as the state space of M, for the specific purpose of the “byte code verification”, i.e., of the verification of safety properties of byte code programs.

The method of operation of the model checker is explained in the following. For each condition s from conditional set S, the model checker verifies if it is fulfilled by state transition system M. If a condition s is not fulfilled, the model checker generates information that can be used together with the knowledge of the instruction for which s was produced, to analyze how the conditions for a secure execution of the JVM can be violated by the byte code program. A typical output of the model checker is made up of the indication of an execution path, i.e., of a sequence of M-states, which leads to a

violation of certain conditions.

A precise analysis of the execution path is not needed. The aim of the byte code verification is merely to establish whether or not the program can violate the conditions. Information pertaining to how a violation can take place may be interesting, but is not needed, since it is merely necessary to reject potentially unsafe programs. In the process, one has to accept the fact that safe programs will also be rejected, which do, in fact, violate conditions, but, when more thoroughly analyzed, would reveal that in an actual environment, they would not be able to cause damage. Since programs of this kind are not generally developed by compilers, but rather only by deliberate encoding, significant restrictions are not being made in this case.

Controlling the byte code verification lies in coordinating the various components. A class data set is subjected to the byte code verification, in that

1. the byte code programs (methods) are released individually;
2. the references in the byte code program are resolved (pre-processing);
3. a state system and a conditional set are constructed;
4. the model checker is started with this input;
5. given success of the model checker, it is repeated for the next byte code program; in the case of failure, it is reported that the byte code verification for the class data set being analyzed has failed.

The method in accordance with the present invention can still be expanded by a few points. Which properties can be verified depends on the manner in which a byte code program is mapped onto a finite state system. To establish, for instance, the acceptability of a byte code program, one can use the mapping outlined above. By selecting another abstraction mapping, one can verify other properties. An interesting property would be, for instance, restricting to a certain extent the consumption of resources by a byte code program. Through the method of the present invention, the byte code verification concept that is critical to safety can be implemented on a formal basis. In this manner, the greatest possible safety is attained for this aspect of Java

technology. One can expect that, above and beyond that, the technique can also be used to verify more far-reaching properties of Applets, making a larger range of applications interesting.

5 This technique is particularly interesting in the environment of Java-capable smart cards (Java-capable meaning the ability to execute byte code programs) where the safety requirements are especially stringent. Java-capable smart cards make it possible to load and execute programs and, in fact, even when they are already in the possession of the end customer (this is not possible or only possible to a limited extent when
10 conventional smart cards are used). In this context, it is extremely important that only those programs are loaded and executed which do not violate the integrity of the card and the data stored thereon, since misusing such data can cause considerable personal or financial damage.

15 The safety of byte code programs can be guaranteed by using the described technique, and a certain functionality can be assured through additional upgrades. This gives one greater confidence in applications designed to be run on platforms that are critical to safety, such as smart cards.

Reference Symbol List

5	JVM	interpreter for Java byte code programs (Java virtual machine)
	M	finite state transition system
	S	conditional set
10	s	condition of S
	V	instruction directory for byte code instructions
15	C1	conditions that have to be met by a state of the JVM
	C2	conditions that have to be met by the JVM
	A1	property that applies for conditions in V
20	B	program instruction
	P	parameters for B

Patent Claim

1. A method for verifying the safety properties of Java byte code programs in accordance with the principle of byte code verification, characterized in that
 - a) using an algorithm describing the properties of byte code instructions, the method of functioning of the byte code program to be verified is mapped by a potentially infinite state transition system onto a finite state transition system (M), and the state space of the interpreter (JVM) is mapped onto a finite set of states in the finite state transition system (M), all information not needed for the acceptability of the byte code program to be checked being omitted, so that the resulting finite state transition system (M) exclusively contains type information, which is used to verify the acceptability of the byte code program and which is entered into a model checker; that
 - b) the properties which characterize an acceptable byte code program are acquired in a logic operation in the form of formulas and are entered as conditional set (S) into the model checker, the model checker interpreting each individual condition (s) of the conditional set (S) as a specification language for the system properties of byte code programs, and that
for each condition (s) of the conditional set (S), the model checker verifies whether it is fulfilled by the state transition system (M), and that the verified byte code program is automatically released for further processing when the state transition system (M) fulfills all conditions (s) of conditional set (S).

Abstract

2. A method for verifying the safety properties of Java byte code programs

2.1 The technical objective is directed to a method which will ensure a greatest possible security in verifying the safety properties of byte code programs.

2.2 In accordance with the present invention, the method of functioning of the byte code program to be verified is modeled on a finite state transition system M, and the state space of the JVM on a finite set of states in M. Once entered into a model checker, the data of finite state transition system M are compared to the data entered in the model checker as conditional set S to determine properties characterizing an acceptable byte code program. The byte code program to be checked is only released for further processing when the state transition system M fulfills all conditions of conditional set S.

2.3 The safety of byte code programs can be guaranteed by using the described technique, and a certain functionality can be assured through additional upgrades. This gives one greater confidence in applications designed to be run on platforms that are critical to safety, such as smart cards.

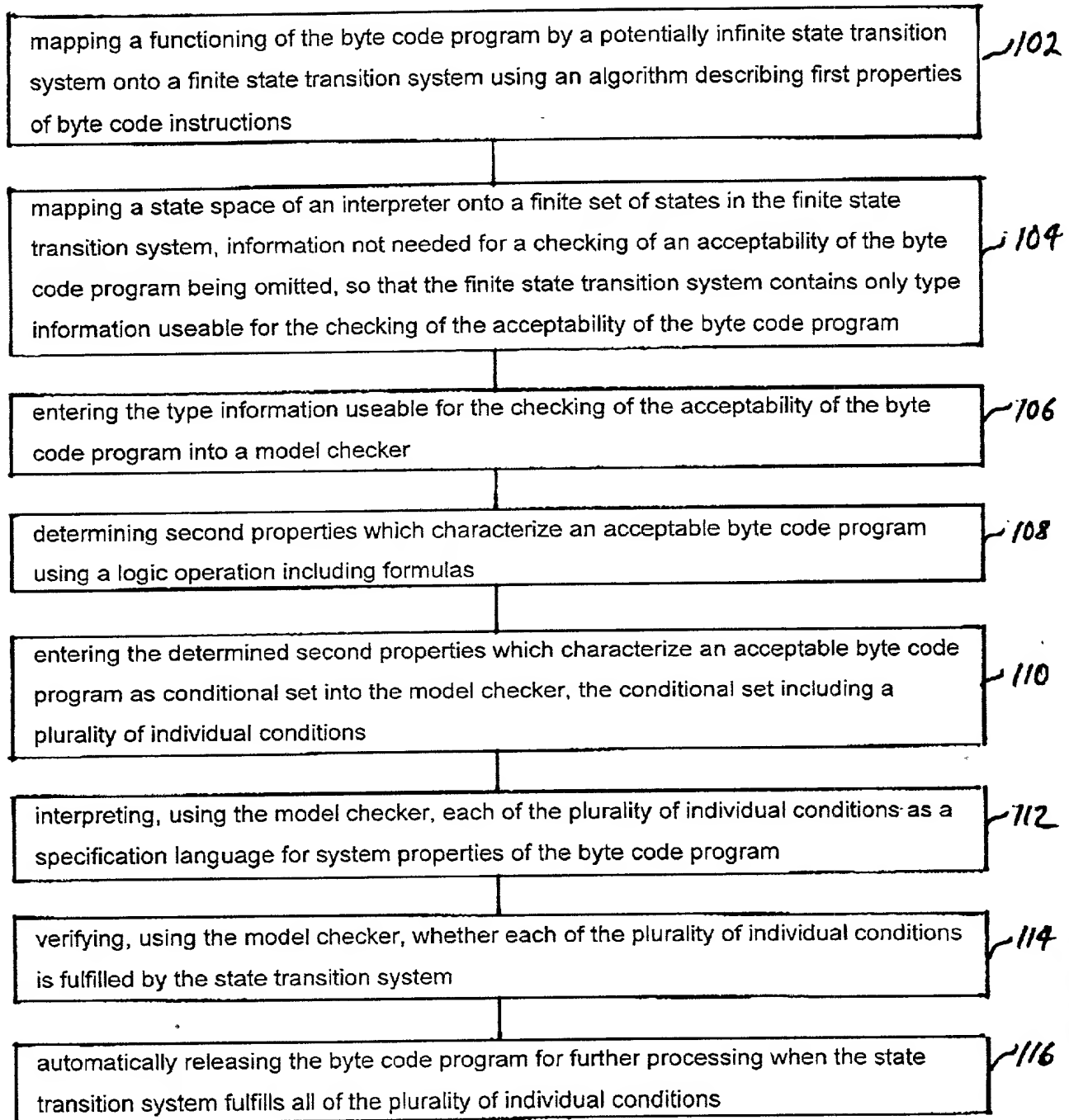


FIG. 1

U.S. DEPARTMENT OF COMMERCE
PATENT AND TRADEMARK OFFICE

DECLARATION AND POWER OF ATTORNEY

ATTORNEY'S DOCKET
NO.
2345/131

As a below named inventor, I hereby declare that:

My residence, post office address, and citizenship are as stated below next to my name,

I believe I am an original, first, and joint inventor of the subject matter that is claimed and for which a patent is sought on the invention entitled **METHOD FOR CHECKING JAVA BYTE CODE PROGRAMMES FOR SECURITY CHARACTERISTICS**, the specification of which was filed as International Application No. **PCT/EP99/04438** on **25 June 1999**.

I hereby state that I have reviewed and understand the contents of the above identified specification, including the claims.

I acknowledge the duty to disclose information which is material to the examination of this application in accordance with Title 37, Code of Federal Regulations, § 1.56(a).

PRIOR FOREIGN APPLICATION(S)

I hereby claim foreign priority benefits under Title 35, United States Code, § 119 of any foreign application(s) for patent or inventor's certificate listed below and have also identified below any foreign application for patent or inventor's certificate having a filing date before that of the application on which priority is claimed:

COUNTRY	APPLICATION NUMBER	DATE OF FILING (day, month, year)	DATE OF ISSUE (day, month, year)	PRIORITY CLAIMED UNDER 35 U.S.C. § 119
Germany	198 30 015.8	26 June 1998		YES

POWER OF ATTORNEY: As a named inventor, I hereby appoint the following attorneys:

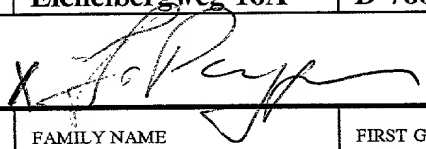
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I declare that all statements made herein of my own knowledge are true and all statements made on information and belief are believed to be true; and further that these statements were made with the knowledge that willful false statements and the like so made are punishable by fine or imprisonment, or both, under § 1001 of Title 18 of the United States Code and that such willful statements may jeopardize the validity of the application or any patent issuing thereon.

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I declare that all statements made herein of my own knowledge are true and all statements made on information and belief are believed to be true; and further that these statements were made with the knowledge that willful false statements and the like so made are punishable by fine or imprisonment, or both, under § 1001 of Title 18 of the United States Code and that such willful statements may jeopardize the validity of the application or any patent issuing thereon.

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Signature X Harald Vegt		Date 21.06.2000	

* An der Dreifaltigkeit 7